

## NAG Library Chapter Introduction

### X02 – Machine Constants

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## 1 Scope of the Chapter

This chapter is concerned with **parameters** which characterise certain aspects of the **computing environment** in which the NAG Library is implemented. They relate primarily to floating-point arithmetic, but also to integer arithmetic, the elementary functions and exception handling. The values of the parameters vary from one implementation of the Library to another, but within the context of a single implementation they are constants.

The parameters are intended for use primarily by other routines in the Library, but users of the Library may sometimes need to refer to them directly.

## 2 Background to the Problems

### 2.1 Floating-point Arithmetic

#### 2.1.1 A model of floating-point arithmetic

In order to characterise the important properties of floating-point arithmetic by means of a small number of parameters, NAG uses a simplified **model** of floating-point arithmetic. The parameters of the model can be chosen to provide a sufficiently close description of the behaviour of actual implementations of floating-point arithmetic, but not, in general, an exact description; actual implementations vary too much in the details of how numbers are represented or arithmetic operations are performed.

The model is based on that developed by Brown (1981), but differs in some respects. The essential features are summarised here.

The model is characterised by four integer parameters. The four integer parameters are:

- $b$ : the base
- $p$ : the precision (i.e., the number of significant base- $b$  digits)
- $e_{\min}$ : the minimum exponent
- $e_{\max}$ : the maximum exponent

These parameters define a set of numerical values of the form:

$$f \times b^e$$

where the exponent  $e$  must lie in the range  $[e_{\min}, e_{\max}]$ , and the fraction  $f$  (also called the mantissa or significand) lies in the range  $[1/b, 1)$ , and may be written

$$f = 0.f_1f_2 \cdots f_p$$

Thus  $f$  is a  $p$ -digit fraction to the base  $b$ ; the  $f_i$  are the base- $b$  digits of the fraction: they are integers in the range 0 to  $b - 1$ , and the leading digit  $f_1$  must not be zero.

The set of values so defined (together with zero) are called **model numbers**. For example, if  $b = 10$ ,  $p = 5$ ,  $e_{\min} = -99$  and  $e_{\max} = +99$ , then a typical model number is  $0.12345 \times 10^{67}$ .

The model numbers must obey certain rules for the computed results of the following basic arithmetic operations: addition, subtraction, multiplication, negation, absolute value, and comparisons: the computed result must be the nearest model number to the exact result (assuming that overflow or underflow does not occur); if the exact result is midway between two model numbers, then it may be rounded either way.

For division and square root, this latter rule is relaxed: the computed result may also be one of the next adjacent model numbers on either side of the permitted values just stated.

On many machines, the full set of representable floating-point numbers conforms to the rules of the model with appropriate values of  $b$ ,  $p$ ,  $e_{\min}$  and  $e_{\max}$ . For machines supporting IEEE binary double precision arithmetic:

$$\begin{aligned}
 b &= 2 \\
 p &= 53 \\
 e_{\min} &= -1021 \\
 e_{\max} &= 1024.
 \end{aligned}$$

(**Note:** the model used here differs from that described in Brown (1981) in the following respect: square-root is treated, like division, as a weakly supported operator.)

### 2.1.2 Derived parameters of floating-point arithmetic

Most numerical algorithms require access, not to the basic parameters of the model, but to certain derived values, of which the most important are:

$$\begin{aligned}
 \text{the machine precision } \epsilon: &= \left(\frac{1}{2}\right) \times b^{1-p} \\
 \text{the smallest positive model number:} &= b^{e_{\min} - 1} \\
 \text{the largest positive model number:} &= (1 - b^{-p}) \times b^{e_{\max}}
 \end{aligned}$$

It is important to note that the machine precision defined here differs from that defined by ISO Fortran 95 (1997).

Two additional derived values are used in the NAG Library. Their definitions depend not only on the properties of the basic arithmetic operations just considered, but also on properties of some of the elementary functions. We define the **safe range** parameter to be the smallest positive model number  $z$  such that for any  $x$  in the range  $[z, 1/z]$  the following can be computed without undue loss of accuracy, overflow, underflow or other error:

$$\begin{aligned}
 &-x \\
 &1/x \\
 &-1/x \\
 &\sqrt{x} \\
 &\log(x) \\
 &\exp(\log(x)) \\
 &y^{(\log(x)/\log(y))} \text{ for any } y
 \end{aligned}$$

In a similar fashion we define the safe range parameter for complex arithmetic as the smallest positive model number  $z$  such that for any  $x$  in the range  $[z, 1/z]$  the following can be computed without any undue loss of accuracy, overflow, underflow or other error:

$$\begin{aligned}
 &-w \\
 &1/w \\
 &-1/w \\
 &\sqrt{w} \\
 &\log(w) \\
 &\exp(\log(w)) \\
 &y^{(\log(w)/\log(y))} \text{ for any } y \\
 &|w|
 \end{aligned}$$

where  $w$  is any of  $x$ ,  $ix$ ,  $x + ix$ ,  $1/x$ ,  $i/x$ ,  $1/x + i/x$ , and  $i$  is the square root of  $-1$ .

This parameter was introduced to take account of the quality of complex arithmetic on the machine. On machines with well implemented complex arithmetic, its value will differ from that of the real safe range parameter by a small multiplying factor less than 10. For poorly implemented complex arithmetic this factor may be larger by many orders of magnitude.

## 2.2 Other Aspects of the Computing Environment

No attempt has been made to characterise comprehensively any other aspects of the computing environment. The other functions in this chapter provide specific information that is occasionally required by routines in the Library.

## 3 Recommendations on Choice and Use of Available Routines

Derived parameters of model of floating-point arithmetic,	
largest positive model number .....	X02ALF
machine precision .....	X02AJF
safe range.....	X02AMF
safe range of complex floating-point arithmetic.....	X02ANF
smallest positive model number.....	X02AKF
Largest permissible argument for SIN and COS.....	X02AHF
Largest representable integer.....	X02BBF
Maximum number of decimal digits that can be represented .....	X02BEF
Parameters of model of floating-point arithmetic,	
$b$ .....	X02BHF
$e_{\max}$ .....	X02BLF
$e_{\min}$ .....	X02BKF
$p$ .....	X02BJF

## 4 Routines Withdrawn or Scheduled for Withdrawal

The following lists all those routines that have been withdrawn since Mark 19 of the Library or are scheduled for withdrawal at one of the next two marks.

Withdrawn Routine	Mark of Withdrawal	Replacement Routine(s)
X02DAF	24	No longer required
X02DJF	24	No longer required

## 5 References

Brown W S (1981) A simple but realistic model of floating-point computation *ACM Trans. Math. Software* 7 445–480

ISO Fortran 95 (1997) ISO Fortran 95 programming language (ISO/IEC 1539–1:1997)

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